

Comparison of the moore's method and dijkstra's algorithm in determining the shortest path between two points on a map

BY GROUP 18

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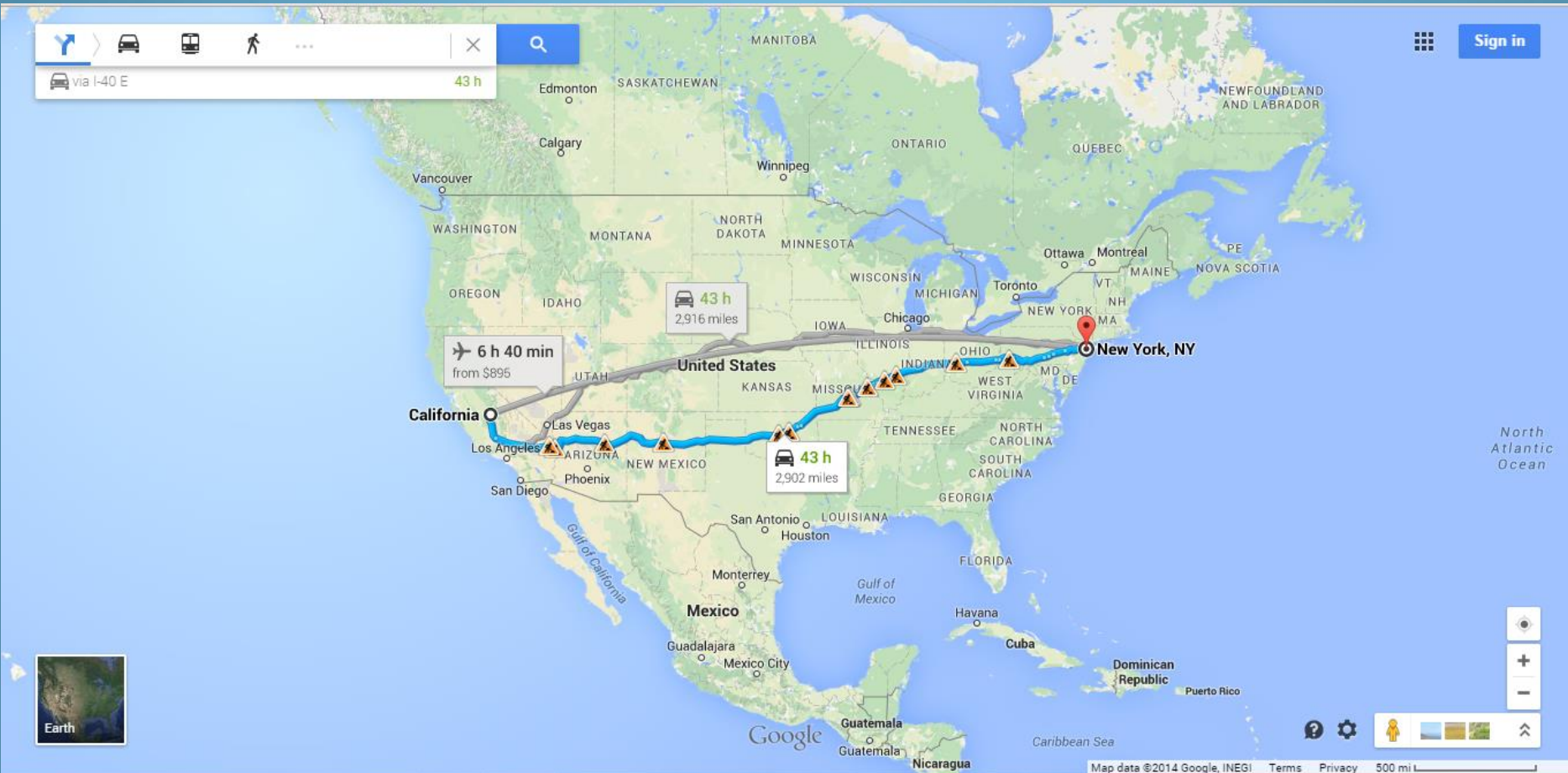
Liu Bowen

Paglia John Naisby

MOTIVATION

- Global Positioning System (GPS) navigation onto smartphone
- Limitations of processing capability
- More efficient algorithm necessary

MOTIVATION

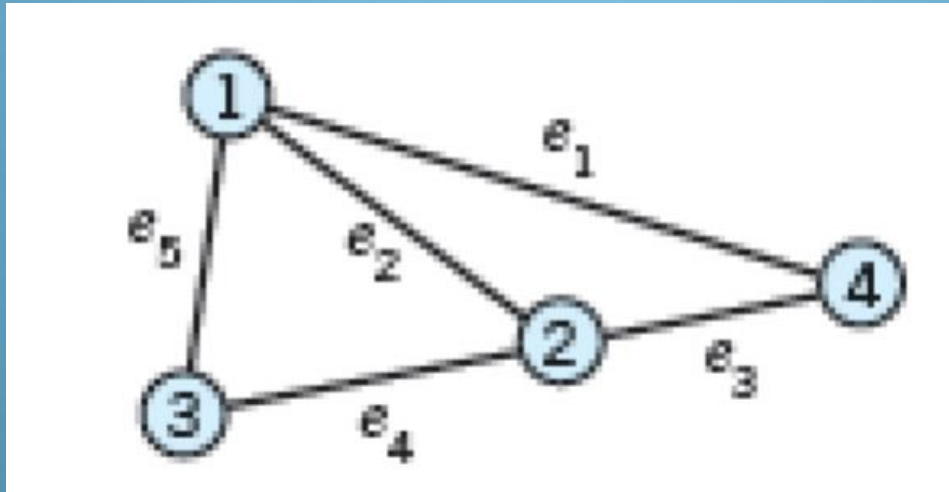


Problem description

- Maps made with $n \times n$ number of tiles
- A starting point and ending point would be determined within this map
- Maps can be represented as graphs
- Compare Moore's and Dijkstras method for these graphs

Introduction to methods

Background Knowledge



- ▷ A set V of vertices: stand for spots in map models
- ▷ A set E of edges: stand for routes connecting spots

Moore's Method (Breadth-First-Search Algorithm)

- ▷ The widely used fundamental method in searching for the shortest path
- ▷ Blind searching method: searching every node of a map to find out the shortest path

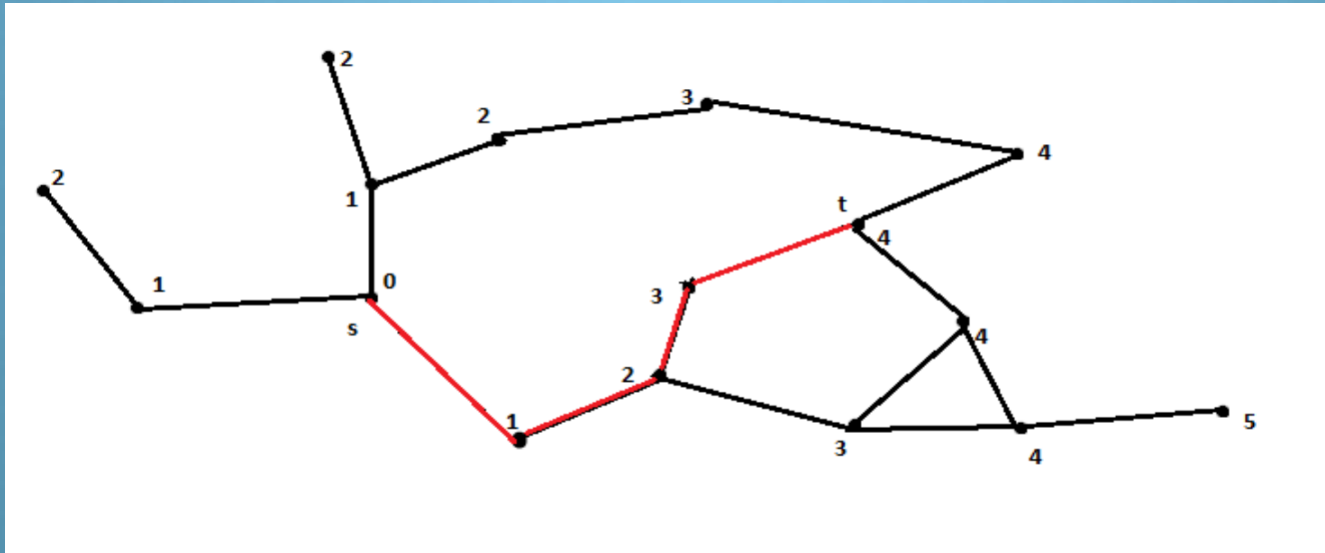
Procedure of the BFS method

1. Label the start point s with 0;
2. Set $i=0$;
3. Search all unlabeled vertices adjacent to a vertex labeled i ;
4. Label all vertices just found with $i+1$;
5. If the terminal point t is found, stop the process and output the shortest path. Otherwise, loop back to the step 3 and keep going on.

IMPORTANT ASSUMPTION

- ▷ Every edge between any two vertices has the same length.
- ▷ Problem changed from finding out the shortest path to finding out the path with the smallest number of edges.

EXAMPLE



Solution:

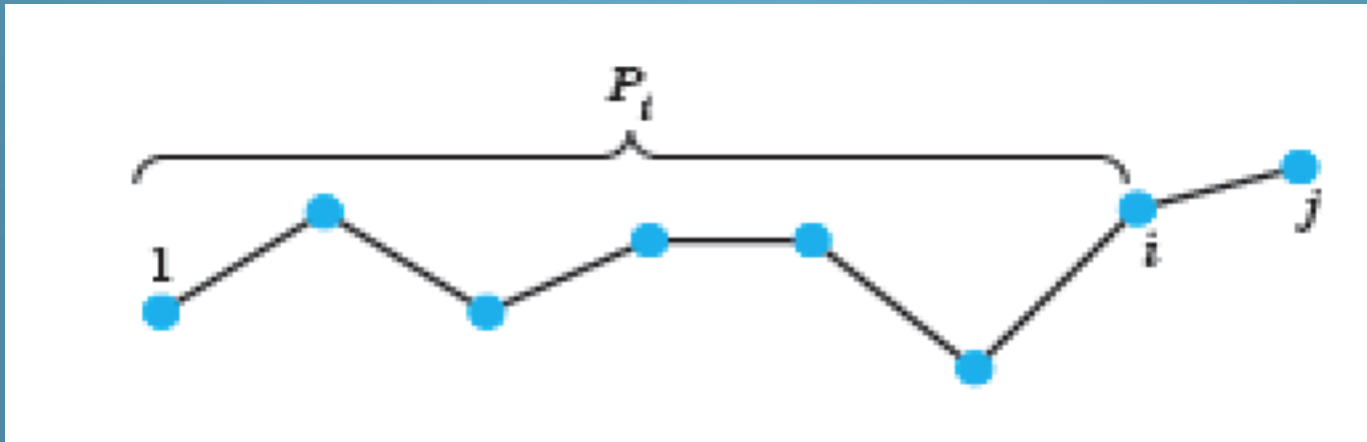
1. Label the start point s with 0;
2. Search all unlabeled vertices adjacent to the point 0 and label them with 1;
3. Repeat the previous works to vertices with new labels, until getting the terminal point t labeled;
4. The number that terminal point labeled with is the shortest length, and the red path is the shortest

Dijkstra's Algorithm

- ▷ Kind of similar to the BFS method
- ▷ Each vertex is labeled in either two ways
 1. A permanent label L_j for the length of a shortest path $1 \rightarrow j$;
 2. A temporary label with stored L_j^* for the length of a shortest path $1 \rightarrow j$;
 3. Set up two sets (stacks) of ps and ts respectively for the permanent labels and temporary ones.

Bellman's Optimality Principle

If $P: 1 \rightarrow j$ is a shortest path from 1 to j in graph G and (i, j) is the last edge of P_j (as picture following), then $P_i: 1 \rightarrow j$ [obtained by dropping (i, j) from P_j] is a shortest path $1 \rightarrow j$.



Pseudocode

Input: Number of vertices n , edges (i, j) , and length l_{ij}

Output: Lengths L_j of the shortest path $1 \rightarrow j$, $j = 2, 3, \dots, n$

1. Initial Step

Vertex 1 gets PL: $L_1=0$

Vertex $j (= 2, \dots, n)$ gets TL: $L_j^*=l_{1j}$ ($=\infty$ if there is no edge (i, j) in G).

Set $ps=\{1\}$, $ts=\{2,3,\dots,n\}$.

2. Fixing a permanent label

Find a k in ts for which L_k^* is minimum, set $L_k = L_k^*$. Take the smallest k if there are several. Delete k from ts and include it in ps .

If $ts = \emptyset$ (that is, ts is empty) then

OUTPUT L_2, \dots, L_n , stop

Else continue (that is, go to Step 3).

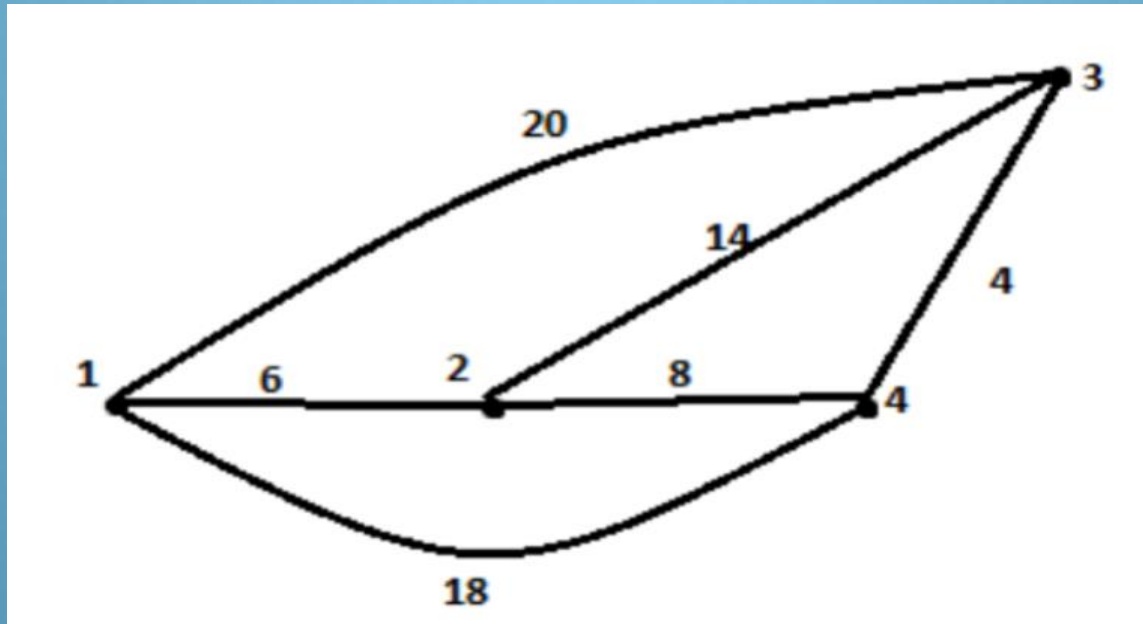
3. Updating temporary labels

For all j in ts , set $L_j^* = \min \{ L_j^*, L_k + l_{kj} \}$

Go to Step 2

END

Example



Question: find out the shortest path 1→2, 3, 4

Solution. The steps and computations.

Step 1. $L_1=0, L_2^*=6, L_3^*=20, L_4^*=18$

Step 2. $L_2 = \min \{L_2^*, L_3^*, L_4^*\} = 6, k=2$

Step 3. $L_3^* = \min \{20, L_2 + L_{23}\} = \min \{20, 6+14\} = 20,$

$L_4^* = \min \{18, L_2 + L_{24}\} = \min \{18, 6+8\} = 14,$

Step 2. $L_4 = \min \{L_3^*, L_4^*\} = 14, k=4$

Step 3. $L_3^* = \min \{20, L_4 + L_{43}\} = \min \{20, 14+4\} = 18$

Step 2. $L_3 = 18, k=3$

$ps = \{1\}, ts = \{2, 3, 4\}$

$ps = \{1,2\}, ts = \{3, 4\}$

$ps = \{1,2,4\}, ts = \{4\}$

$ps = \{1,2,3,4\}, ts = \emptyset$

Figure 1.4 shows the resulting shortest paths, of lengths $L_2=6$, $L_3=18$, $L_4=14$

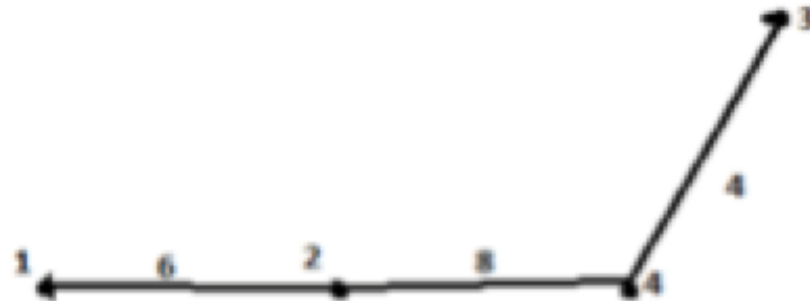


Figure 1.4 - Shortest Paths From Vertex 1 To Vertices 2, 3, 4

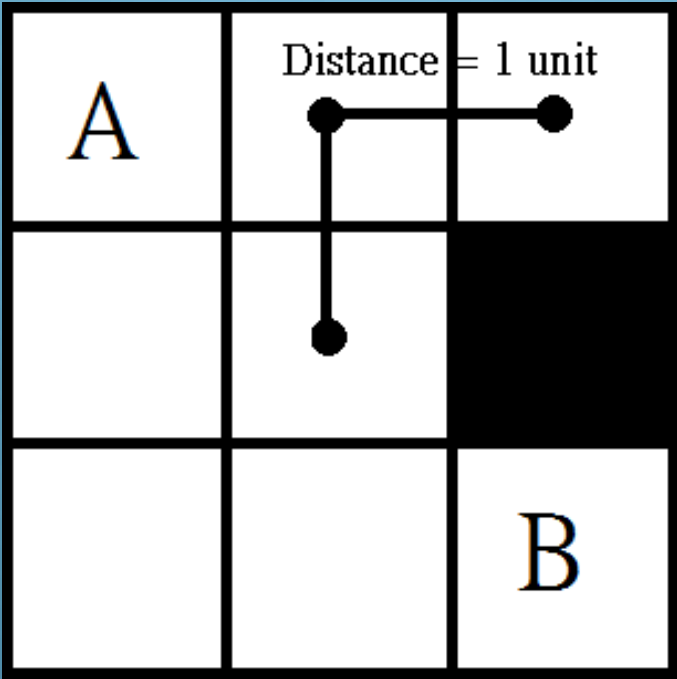
Mathematical presentation of problem

Assumptions for this problem

- The distance between the nodes is 1 unit.
- Turning a corner has no impact on the distance travelled.
- The edge between nodes are bidirectional

Mathematical presentation of problem

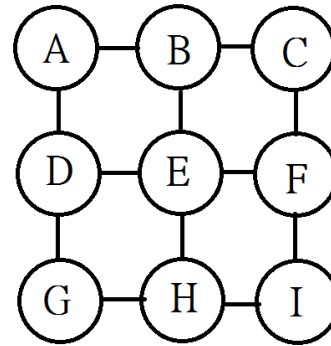
1. create maps



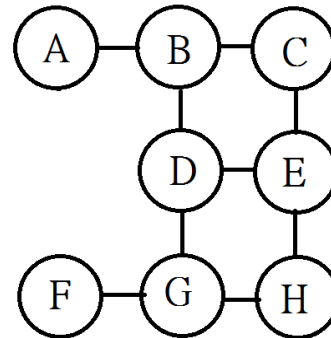
Mathematical presentation of problem

2. transfer maps into graphs

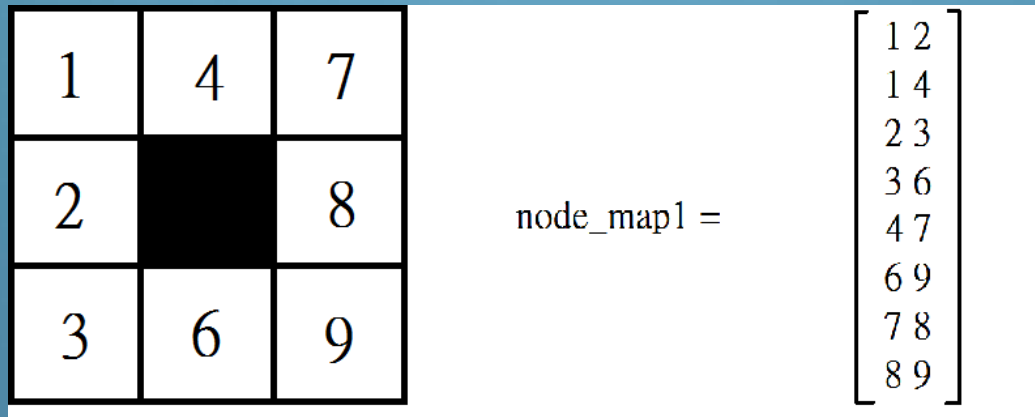
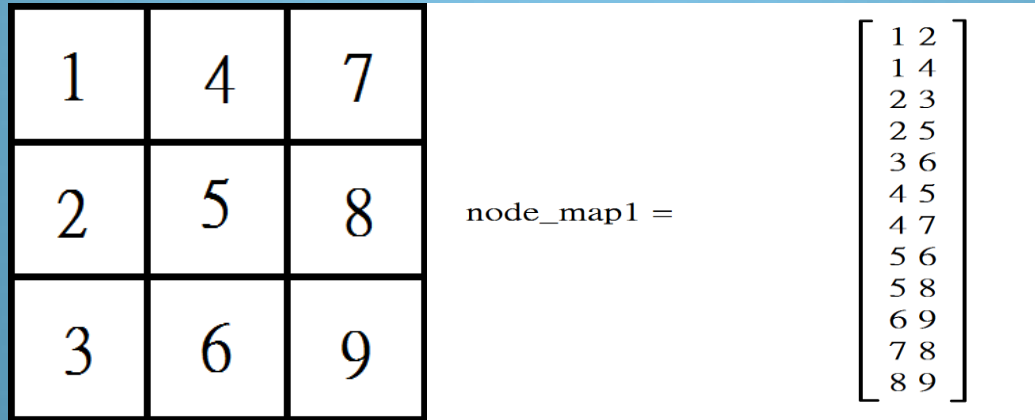
A	B	C
D	E	F
G	H	I



A	B	C
	D	E
F	G	H

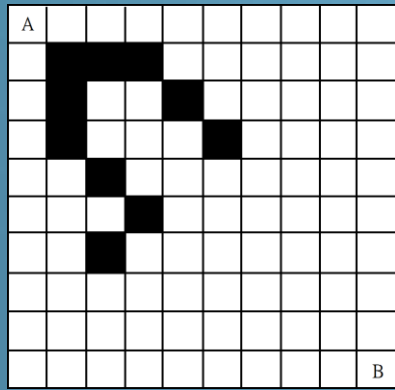


Transforming the designed maps into graphs

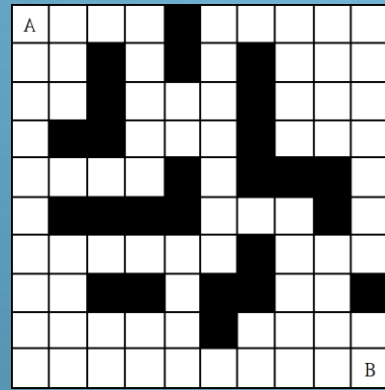


Mathematical presentation of problem

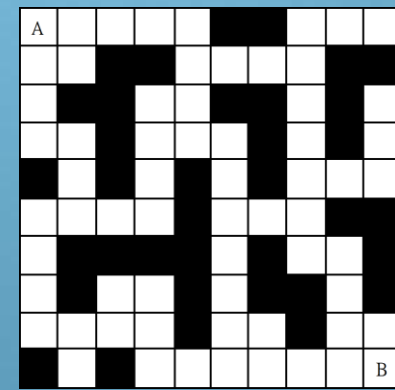
maps



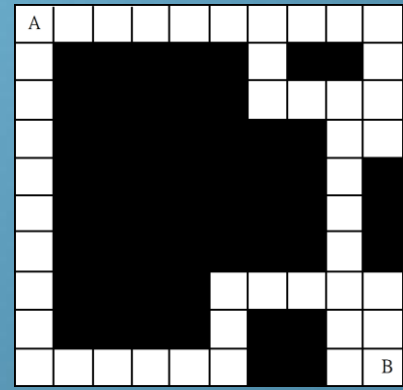
(a)



(b)



(c)

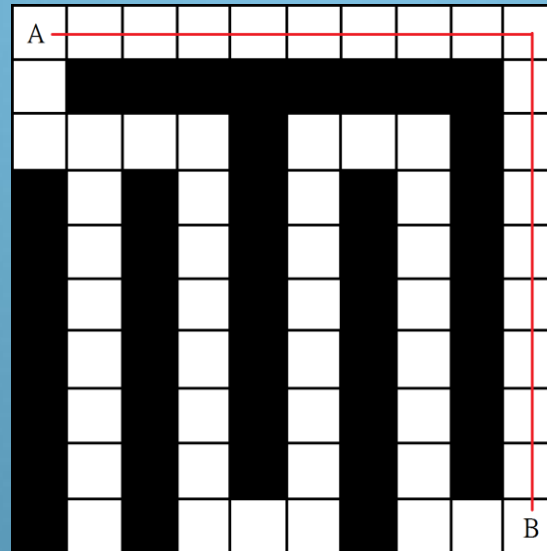


(d)

Maps (a), (b), (c), (d) have been developed with 10, 25, 36 and 55 obstacles

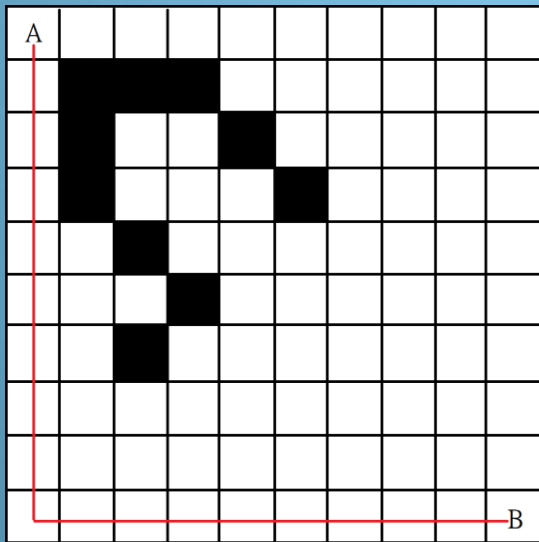
Reference map

Upon the application of the MATLAB functions to the reference map, the same shortest path was found successfully by both methods.

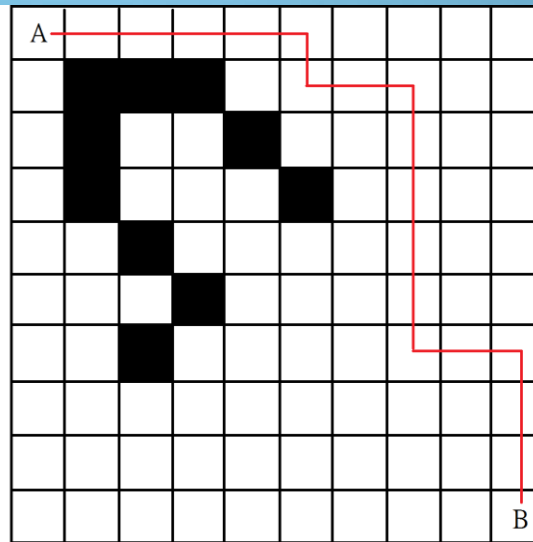


	Moore's	Dijkstra's
Reference Map (average time)	0.000733s	0.000743s
Shortest path (node – node)	1-11-21-31-41-51-61- 71-81-91-92-93-94-95- 96-97-98-99- 100(18steps)	1-11-21-31-41-51-61- 71-81-91-92-93-94-95- 96-97-98-99- 100(18steps)

Maps with different number of obstacles



(a)

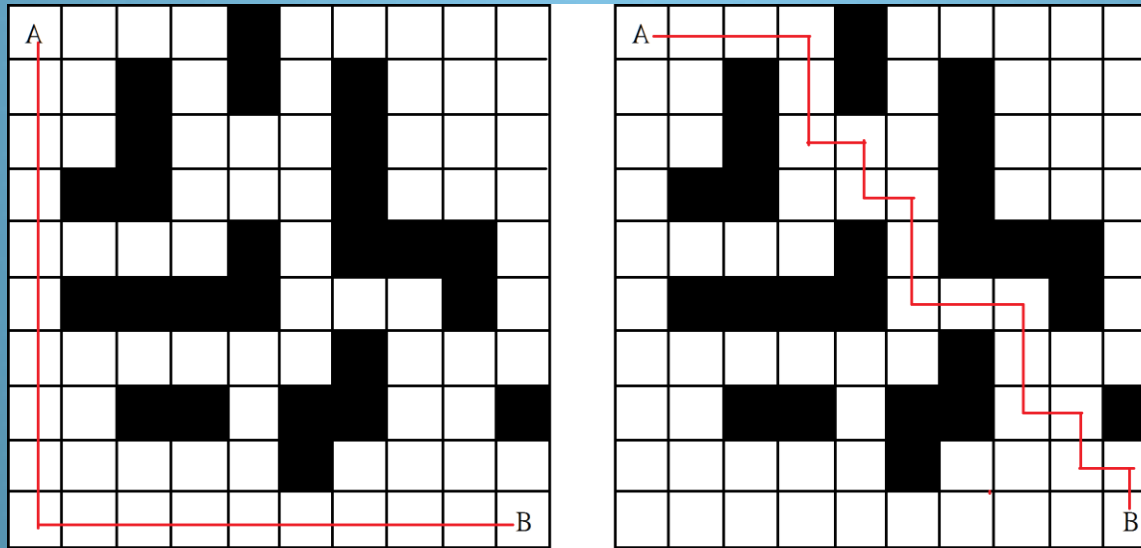


(b)

(a) Moore's method (b) Dijkstra's Algorithm

	Moore's	Dijkstra's
Map10 (average time)	0.000813s	0.000848s
Shortest path (node – node)	1- 2- 3- 4- 5- 6- 7- 8- 9-10-20-30-40-50-80-90-100(18 steps)	1-11-21-31-41-51-52-62-72-73-74-75-76-77-87-97- 98-99-100 (18 steps)

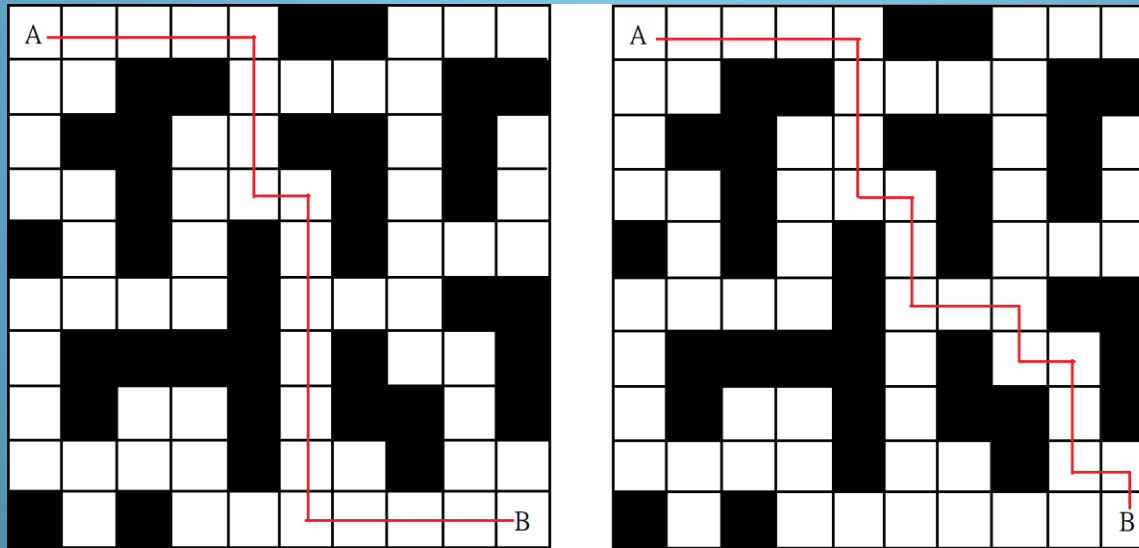
Maps with different number of obstacles



(a) Moore's method (b) Dijkstra's Algorithm

	Moore's	Dijkstra's
Map25 (average time)	0.000750s	0.000781s
Shortest path (node – node)	1- 2- 3- 4- 5- 6- 7- 8- 9-10-20-30-40-50-80-90-100(18 steps)	1-11-21-31-32-33-43-44-54-55-56-66-76-77-78-88-89-99-100 (18 steps)

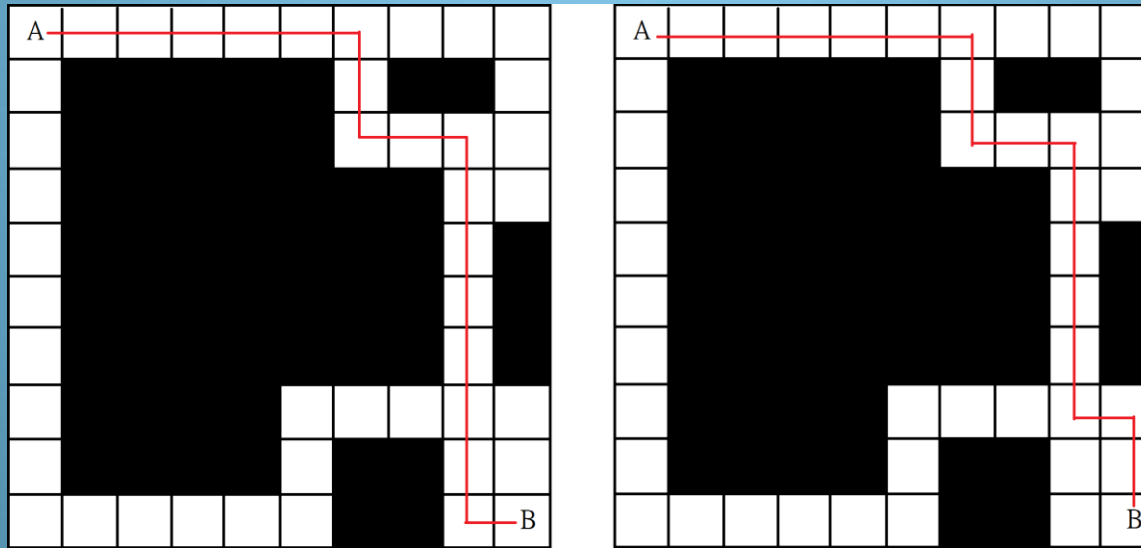
Maps with different number of obstacles



(a) Moore's method (b) Dijkstra's Algorithm

	Moore's	Dijkstra's
Map36 (average time)	0.000724s	0.000769s
Shortest path (node – node)	1-11-21-31-41-42-43-44-54-55-56-57-58-59-60-70-80-90-100(18 steps)	1-11-21-31-41-42-43-44-54-55-56-66-76-77-87-88-89-99-100(18 steps)

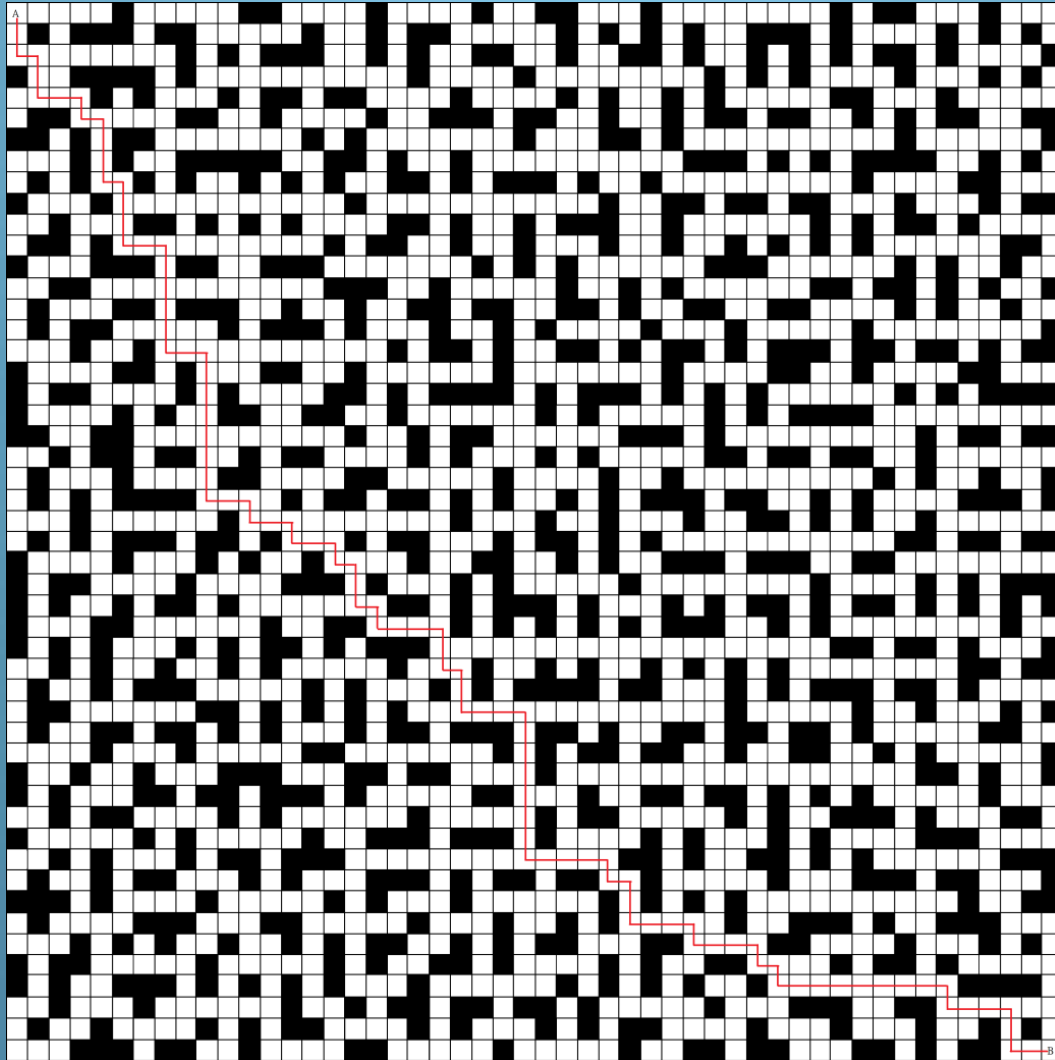
Maps with different number of obstacles



(a) Moore's method (b) Dijkstra's Algorithm

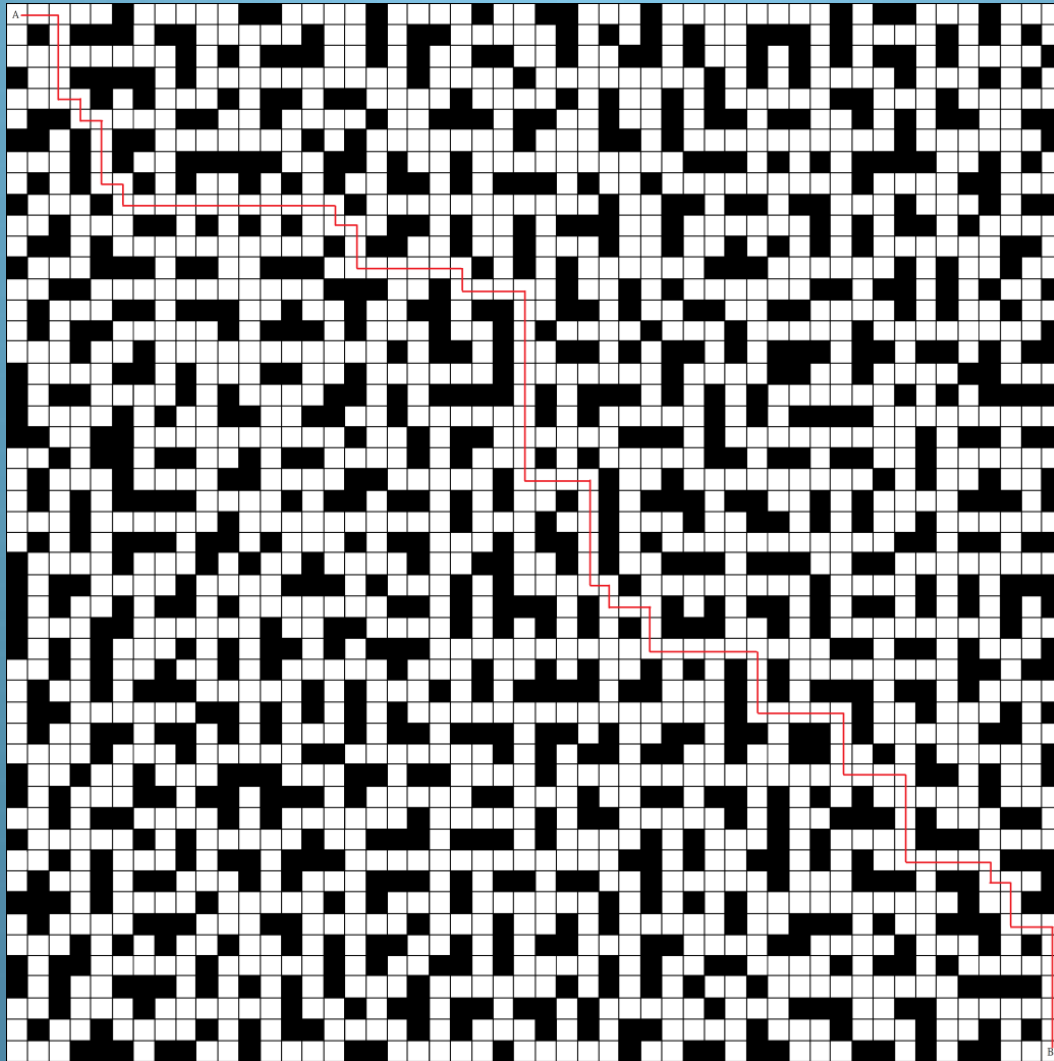
	Moore's	Dijkstra's
Map50 (average time)	0.000713s	0.000742s
Shortest path (node – node)	1-11-21-31-41-51-61-62-63-73-83-84-85-86-87-88-89-90-100(18 steps)	1-11-21-31-41-51-61-62-63-73-83-84-85-86-87-88-98-99-100(18steps)

50x50 Map with 900 obstacles



the shortest path – Moore's method

50x50 Map with 900 obstacles



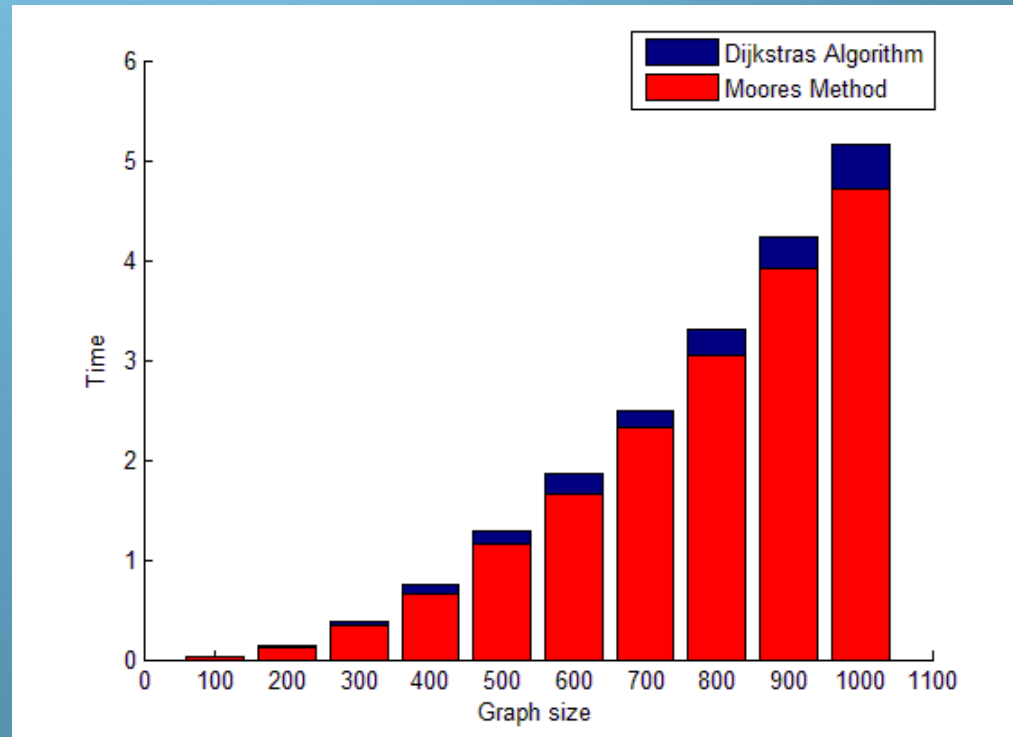
the shortest path – Moore's method

50x50 Map with 900 obstacles

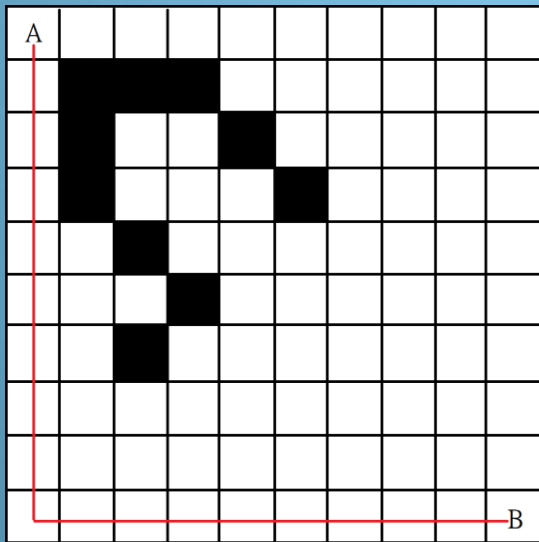
	Moore's	Dijkstra's
Map 50x50 (no obstacles)	0.006550s	0.006604s
Map 50x50 (average time)	0.00444s	0.004850s
Shortest path (node – node)	1-2-3-53-54-55-105-155-156-206-207-208-209-259-260-261-262-312-362-363-364-365-366-367-417-467-468-469-470-471-472-473-474-524-574-575-625-675-676-726-776-777-827-828-829-879-880-930-980-1030-1031-1032-1082-1083-1084-1134-1184-1234-1235-1236-1237-1238-1239-1240-1241-1291-1341-1391-1441-1442-1492-1493-1494-1544-1594-1644-1645-1695-1745-1795-1796-1846-1847-1897-1947-1997-2047-2097-2147-2197-2247-2248-2298-2348-2398-2399-2400-2450-2500 (98 steps)	1-51-101-102-103-104-105-155-156-206-207-208-209-259-260-310-360-410-460-510-560-610-660-710-760-761-811-812-813-863-913-963-1013-1063-1064-1114-1164-1214-1215-1216-1217-1218-1219-1220-1221-1222-1223-1273-1323-1373-1374-1375-1376-1377-1378-1428-1429-1479-1529-1530-1531-1581-1631-1681-1731-1781-1782-1783-1784-1834-1884-1934-1984-1985-1986-1987-2037-2087-2137-2138-2139-2140-2141-2191-2241-2291-2341-2342-2392-2393-2394-2444-2494-2495-2496-2497-2498-2499-2500 (98 steps)

Timing Results

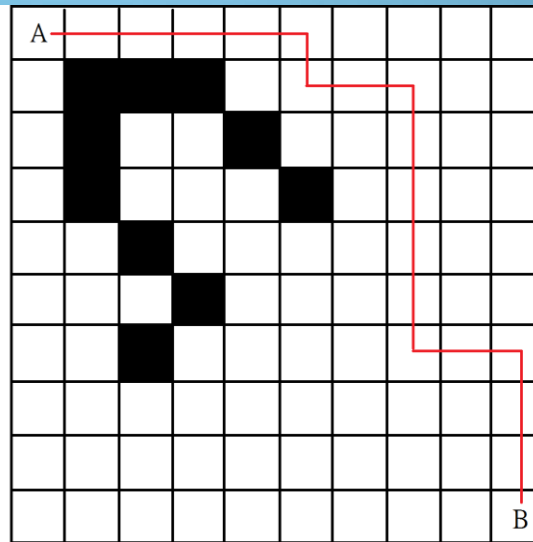
	Moore's	Dijkstra's
100x100	0.023097	0.02613
200x200	0.127722	0.141696
300x300	0.346553	0.383323
400x400	0.661585	0.756057
500x500	1.160261	1.287945
600x600	1.660201	1.852375
700x700	2.321567	2.486482
800x800	3.039056	3.296351
900x900	3.923256	4.225479
1000x1000	4.714555	5.157708



Maps with different number of obstacles



(a)

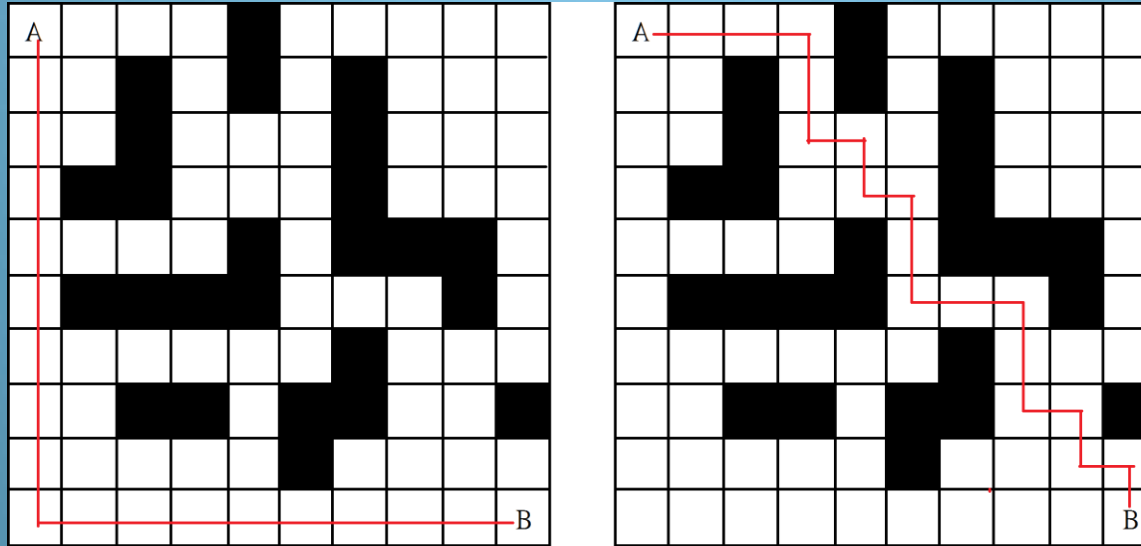


(b)

(a) Moore's method (b) Dijkstra's Algorithm

	Moore's	Dijkstra's
Map10 (average time)	0.000813s	0.000848s
Shortest path (node – node)	1- 2- 3- 4- 5- 6- 7- 8- 9-10-20-30-40-50-80-90-100(18 steps)	1-11-21-31-41-51-52-62-72-73-74-75-76-77-87-97- 98-99-100 (18 steps)

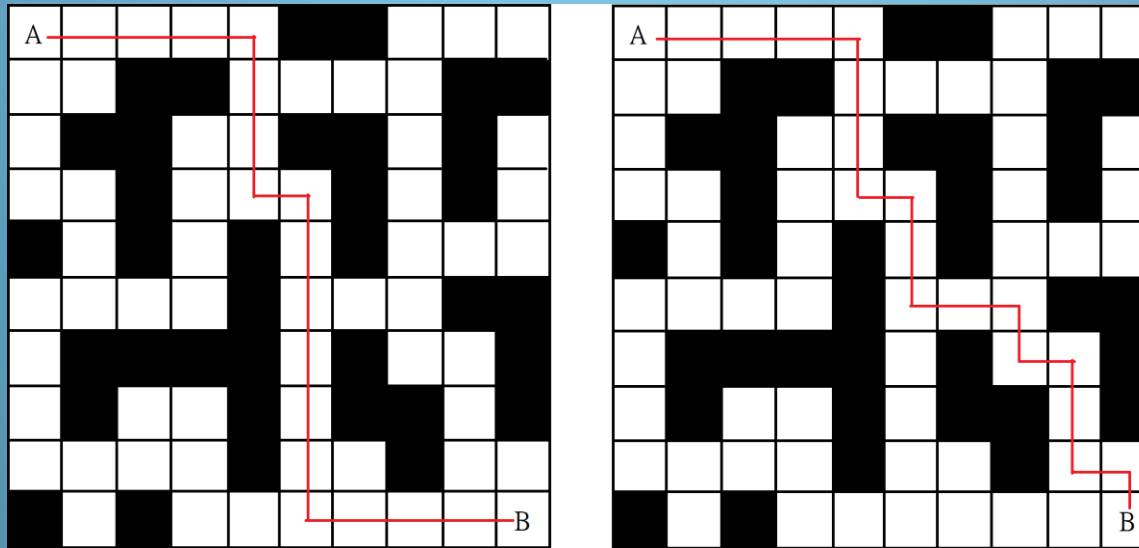
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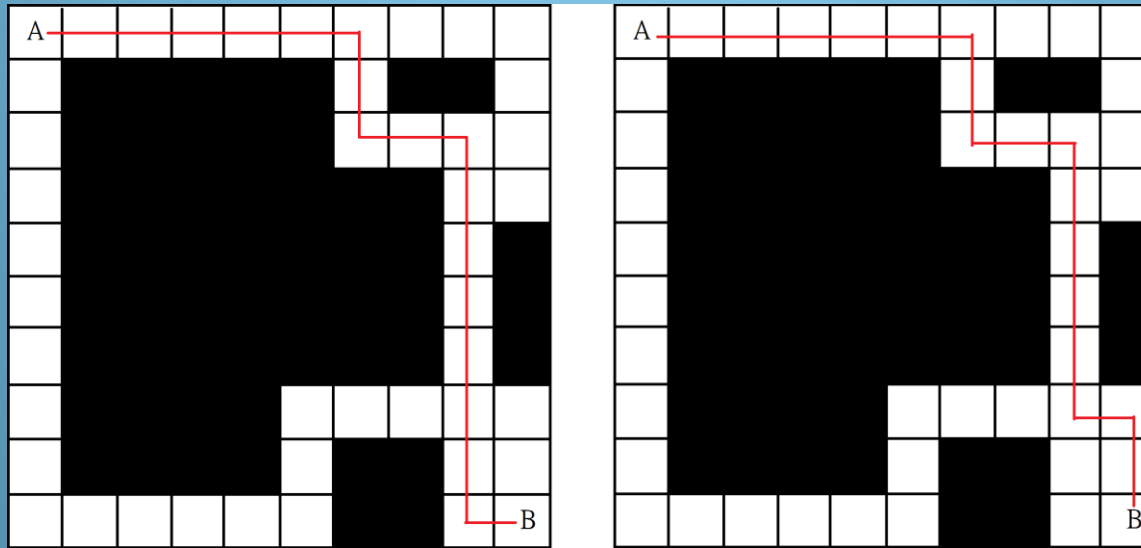
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Maps with different number of obstacles

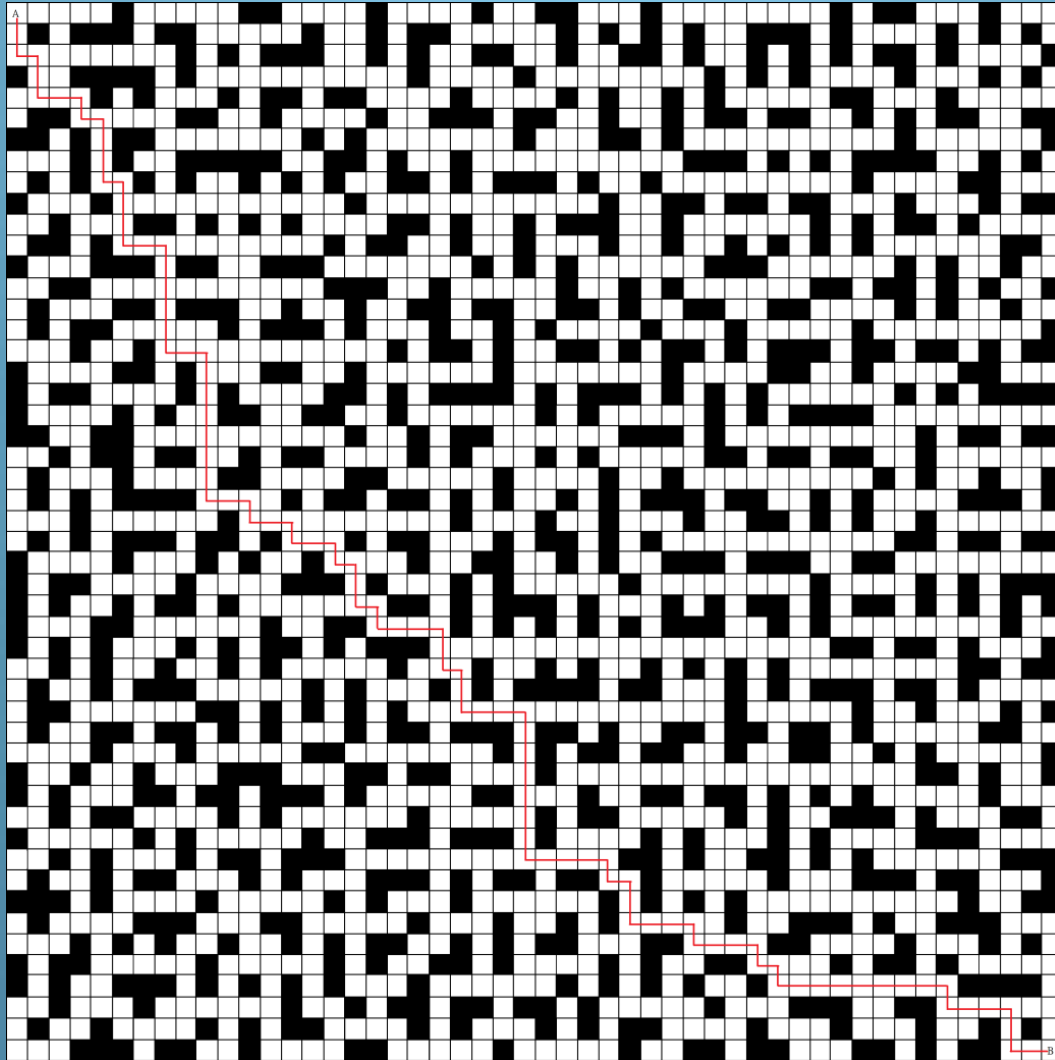


(a)

(b)

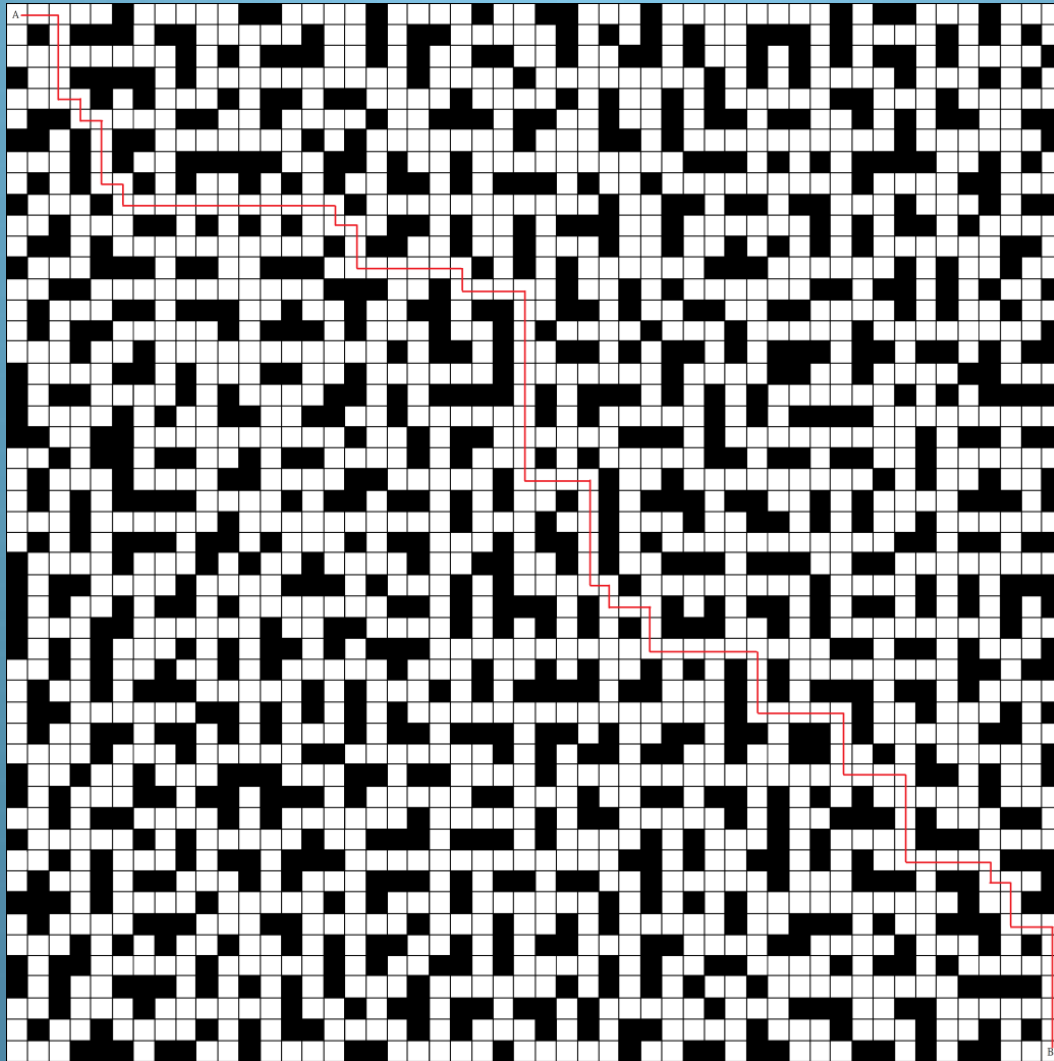
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50x50 Map with 900 obstacles



the shortest path – Moore's method

50x50 Map with 900 obstacles



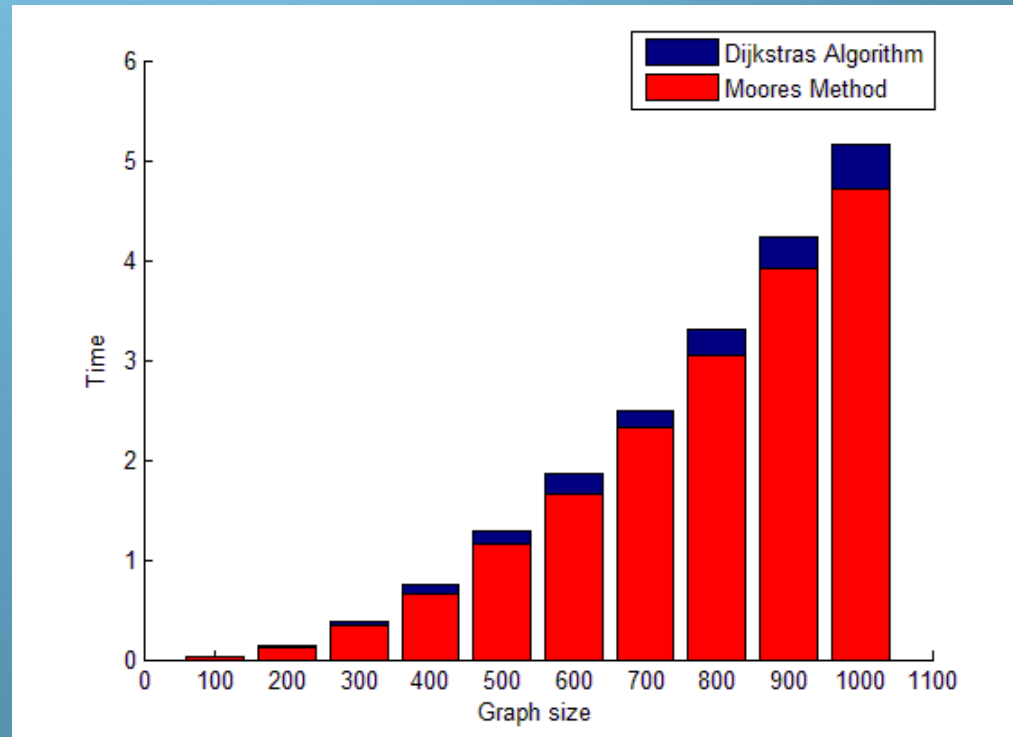
the shortest path – Moore's method

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	Moore's	Dijkstra's
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1000x1000	4.714555	5.157708



Conclusion

- The Moore's method is more time efficient than the Dijkstra's algorithm for the applications in this project
- As the map gets larger, the time required for both methods take longer
- The Dijkstra's algorithm would be more suitable for analysing weighted graphs

Future work

- Improve MATLAB algorithm to produce weighted maps for Dijkstra's algorithm
- Investigate the effect of adding distance weighting to corners